EXTREMAL CLIQUE COVERINGS OF COMPLEMENTARY GRAPHS

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Let cc(G) (resp. cp(G)) be the least number of complete subgraphs needed to cover (resp. partition) the edges of a graph G. We present bounds on $\max \{cc(G) + cc(\bar{G})\}$, $\max \{cp(G) + cp(\bar{G})\}$, $\max \{cp(G)cp(\bar{G})\}$ where the $\max \max$ are taken over all graphs G on n vertices and \bar{G} is the complement of G in K_n . Several related open problems are also given.

Introduction

Let G be a graph on n vertices and let \overline{G} be its complement in K_n , the complete graph on n vertices. If f is a real valued function defined on graphs, what are the extreme values of $f(G)+f(\overline{G})$ and f(G) $f(\overline{G})$? E. A. Nordhaus and J. W. Gaddum (see e.g. [5]) considered those questions when the function is the chromatic number. D. Taylor, R. D. Dutton and R. C. Brigham [5] studied the questions for several other functions. One of those is the clique covering number. That is cc(G), the least number of complete subgraphs (cliques) of G necessary to cover the edge set of G. We continue their investigation. We also consider the questions for another function the clique partition number. That is cp(G), the least number of cliques needed to partition the edge set of G.

In Theorem 1, we establish the right inequality of $\lfloor n^2/4 \rfloor + 2 \le \max \{cc(G) + +cc(\overline{G})\} \le (n^2/4)(1+o(1))$ where the maximum is taken over all graphs G on n vertices. The bipartite graph $K_{\lfloor n/2 \rfloor, \lceil n/2 \rfloor}$ assumes the lower bound.

vertices. The bipartite graph $K_{\lfloor n/2\rfloor,\lceil n/2\rceil}$ assumes the lower bound.

In Theorem 2 we modify the proof of Theorem 1 to show that $\max \{cc(G)cc(\overline{G})\} \le (n^4/256)(1+o(1))$, where the maximum is taken over all n-vertex graphs G. D. Taylor et al. [5, Theorem 5] gave an example of a graph F for which $cc(F)cc(\overline{F})=n^2(n+8)^2/256$. The graph F is obtained from two copies A_1 and A_2 of $K_{n/4}$ and two copies A_2 and A_3 of $\overline{K}_{n/4}$ by joining each vertex of A_i to each vertex of A_{i+1} (i=1,2 and 3). When n is not divisible by 4 the construction can be modified to yield a similar graph. Hence Theorem 2 establishes the conjecture made in [5], that $\max \{cc(G)cc(\overline{G})\} \sim n^4/256$ where the maximum is taken over all n-vertex graphs G.

Somewhat weaker results for the clique partition number are obtained in Theorems 3, 4 and 5. They imply

$$\frac{7}{25} n^2 + O(n) \le \max \{ cp(G) + cp(\overline{G}) \} \le \frac{13}{30} n^2 + O(n) \quad \text{and} \quad$$

$$\frac{29}{2000} n^4 + O(n^3) \le \max\{cp(G)cp(\overline{G})\} \le \frac{169}{3600} n^4 + O(n^3)$$

where the maxima are taken over all n-vertex graphs G.

We state several related open problems at the end of the paper.

Results

Theorem 1. For some d>0 and all graphs G on n vertices, $cc(G)+cc(\overline{G})<<(n^2/4)(1+d/\log n)$.

Proof. Suppose $4^{\varepsilon} \leq n/4c^3$. From a sequence $\mathscr{S} = \{K^1, K^2, ..., K^l\}$ of cliques K^j in K_n by choosing K^i to be a clique in G or in \overline{G} which covers at least c edges uncovered by $K^1 \cup K^2 \cup ... \cup K^{i-1}$. The process halts when such a selection is no longer possible. Now $l \leq n^2/c$. If a vertex has fewer than n/c incident edges in G or in \overline{G} , augment $\mathscr S$ by adding these edges separately, and continue repeating this step until there are no such vertices remaining. At most $2n^2/c$ new cliques have been added to $\mathscr S$. Let H_1 (or H_2) denote the subgraph of K_n induced by the set of edges of G (respectively \overline{G}) not contained in the union of the cliques in $\mathscr S$, and put $H = H_1 \cup H_2$. Let T denote the set of vertices of H with degree at least n/c in both H_1 and H_2 , and let U and V denote the sets of vertices in $K_n - T$ with degree at least n/c in H_1 and H_2 respectively. Note that vertices in U and V have degree 0 in H_2 and H_1 respectively.

In [2] it is shown that $cc(D) \le k^2/4$ for all k-vertex graphs D. Therefore the edges of H with both ends in U or both ends in V can be covered by at most $|U|^2/4$ or $|V|^2/4$ cliques respectively. We further augment $\mathscr S$ by these cliques, which adds at most $n^2/4$ cliques to $\mathscr S$.

We next show that $|T| \le n/c$. Assume |T| > n/c. Then at least $n^2/2c^2$ edges of H_1 have at least one end in T. It follows that some set E of at least $n/2c^2$ such edges are all incident with some vertex p. Let $T' = \{v \in T : pv \in E\}$. Then $|T'| \ge n/2c^2$, so at least $n^2/4c^3$ edges of H_2 have at least one end in T'. Then a set F of $n/4c^3$ or more such edges are all incident with some vertex q. Let $T'' = \{v \in T' : qv \in F\}$. Then $|T''| \ge n/4c^3$. By the bound for Ramsey's Theorem given for example in [1, Theorem 7.5], G of \overline{G} contains a clique K with c vertices in T''. Therefore the clique spanned by K and p (or K and q) covers c edges of H_1 (respectively H_2) incident with p (respectively q). But this contradicts the definition of \mathcal{S} . Thus $|T| \le n/c$ as claimed. Hence we can further augment the cliques in \mathcal{S} by adding all edges of H incident with vertices in T as separate cliques. There are at most n^2/c such edges.

The cliques in \mathcal{S} now form a clique covering of G and a clique covering of G, and $|\mathcal{S}| \le n^2/4 + 4n^2/c$. For large n we can take $3c > \log n$, which gives the theorem.

Theorem 2. For some d>0 and all graphs G on n vertices, $cc(G) \cdot cc(G) < n^4(1+d/\log n)/256$.

Proof. In the proof of Theorem 1, we obtained a clique covering of G using at most $4n^2/c + |U|^2/4$ cliques, and a clique covering of \overline{G} using at most $4n^2/c + |V|^2/4$ cliques, where $4^c \le n/4c^3$ and $|U| + |V| \le n$. Hence $cc(G)cc(\overline{G}) \le (4n^2/c + a^2/4)(4n^2/c + b^2/4)$ where each of these factors is at most $n^2/2$, and $a+b \le n$. This product is at most $4n^4/c + a^2b^2/16$, which is maximised when a=b=n/2. Hence $cc(G)cc(\overline{G}) \le \le 4n^4/c + n^4/256$. Taking $3c > \log n$ as in Theorem 1, we obtain the result.

Corollary. For each graph G on n vertices min $(cc(G), cc(\overline{G})) \le n^2/16(1+o(1))$.

If G_1 and G_2 are vertex-disjoint graphs, then $G_1 \vee G_2$ is the graph formed from the union of G_1 with G_2 by adding edges joining each vertex of G_1 to each vertex of G_2 .

Lemma 1. [3, Theorem 3]. Let $G = A \vee \overline{K}_q$. If A has p vertices and e edges, and the edge-chromatic number $\chi'(A)$ of A is at most q, then cp(H) = pq - e.

We note that $\chi'(K_m)=m$ or m-1 according as m is odd or even. Therefore for all $m \ge 1$,

$$(1) cp(K_m \vee \overline{K}_m) = m^2 - \binom{m}{2}$$

and

(2)
$$cp(K_{m+r} \vee \overline{K}_{2m}) = 2m(m+r) - {m+r \choose 2}$$
 when $0 \le r \le m$.

Let A and B be replicas of K_m and let H_m be the graph diagrammed in Figure 1. There, as in all figures below, a double line joining two graphs G_1 and G_2 indicates that every vertex in G_1 is adjacent to every vertex in G_2 .

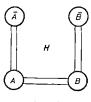


Fig. 1

Lemma 2. For all $m \ge 1$, $cp(H_m) \ge \frac{7}{4} m^2 + m$.

Proof. Let $\mathscr C$ be a clique partition of $H=H_m$ of least cardinality (so that $|\mathscr C|==cp(H)$). Denote the subfamily $\{K^1,K^2,...,K^\sigma\}$, consisting of those cliques in $\mathscr C$ with vertices in both graphs A and B, by $\mathscr S$. From subgraphs A' and B' of A and B by deleting the edges of all cliques in $\mathscr S$ from A and B respectively. Let d_i and e_i be the number of vertices of K^i in A and B respectively. Denote the clique partitions of $\overline{A} \vee A'$ and $\overline{B} \vee B'$ induced by $\mathscr C - \mathscr S$ by $\mathscr C_A$ and $\mathscr C_B$ respectively. Thus $cp(H) = |\mathscr C_A| + |\mathscr C_B| + \sigma$. But

$$|\mathscr{C}_A| \ge cp(\overline{A} \vee A') = m^2 - {m \choose 2} + \sum_{i=1}^{\sigma} {d_i \choose 2}$$

by Lemma 1. Similar statements for B imply that

$$cp(H) \ge m^2 + m + \sigma + \sum_{i=1}^{\sigma} {d_i \choose 2} + {e_i \choose 2}.$$

Differentiation shows that the minimum of the quantities

$$\frac{\binom{d}{2}+\binom{e}{2}+1}{de}$$
,

where d, e are positive integers and $de \ge 1$, is 3/4. This minimum is achieved at d=e=2. Now every edge with one vertex in A and the other in B must be covered by some member of \mathcal{S} . Also K^i in \mathcal{S} covers exactly $d_i e_i$ edges joining A to B. Thus

$$\sum_{i=1}^{\sigma} d_i e_i = m^2$$

and hence $cp(H) \ge 7/4m^2 + m$.

Theorem 3. Let r be the remainder when n is divided by 5. For each $n \ge 20$

$$\max \{cp(G) + cp(\overline{G})\} \ge \frac{7n^2}{25} + \frac{(25+2r)n - 41r^2}{50},$$

where the maximum is over all graphs G on n vertices.

Proof. Let L be a replica of K_{m+r} and let K be a replica of K_m . Define G_n to be the graph whose diagram is given in Figure 2 (a). The diagram of \overline{G}_n is given in Figure 2 (b). (We use the same diagrammatic convention here as for Figure 1.)

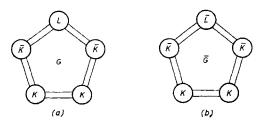


Fig. 2

The graph $G = G_n$ is the edge-disjoint union of $H \equiv H_m$ and $H' \equiv K_{m+r} \vee \overline{K}_{2m}$. Since every clique in G has all its edges in H or all its edges in H', we have cp(G) = cp(H) + cp(H'). Similarly $cp(\overline{G}) = cp(H) + 2cp(\overline{K} \vee \overline{L})$. Since $n \ge 20$, $m \ge 4$ and so equations (1) and (2) imply

(3)
$$\frac{7n^2}{25} + \frac{25n + 2nr - 41r^2}{50} \le cp(G_n) + cp(\overline{G}_n). \quad \blacksquare$$

When does equality hold in (3)? It is a direct consequence of the following Lemma that equality holds infinitely often.

Lemma 3. [4, proof of Theorem 4, pp. 346, 347]. Let K(q, k) be the complete kpartite graph defined by k vertex-disjoint replicas of \overline{K}_q . Then the edge set of K(q, k)can be partitioned into cliques of order k if there exist k-2 mutually orthogonal Latin squares on q symbols.

With k=4, Lemma 3 implies that the edges joining A to B in the graph H_m of Lemma 2 can be covered using edge-disjoint replicas of K_4 for even m>12. Therefore when n>64 and (n-r)/5 is even, equality holds in (3).

Theorem 4. For each graph G on n vertices, $cp(G) + cp(\overline{G}) \le 13n^2/30 - n/6$.

Proof. Let us construct a clique partition of K_n into triangles and edges, each of which is in G or \overline{G} . First select as many edge-disjoint triangles as possible. Then the set of s edges uncovered by any of these t triangles cannot contain the edge set of a copy of K_6 , for otherwise G or \overline{G} would contain a triangle by an instance of Ramsey's theorem. Therefore, by Turán's theorem (see e.g. [1, Theorem 7.9]),

 $s \le 2n^2/5$. Since $3t+s=\binom{n}{2}$, it follows that the partition has at most $13n^2/30-n/6$ members.

The coefficient of n^2 appearing in the right side of the inequality of Theorem 4 can be reduced by 1/204 by using K_4 's as well as K_3 's and K_2 's in the clique partition, and bounds on higher Ramsey numbers lead to further improvements. However, this approach cannot lead to an exact determination of max $\{cp(G)+cp(\overline{G})\}$. The bound in Theorem 3 is probably nearer to the actual value.

Theorem 5. Taking the maximum over all graphs on n vertices,

$$\frac{39}{2000} n^4 + O(n^3) < \max\{cp(G)cp(\overline{G})\} < \frac{169}{3600} n^4 + O(n^3).$$

Proof. The left inequality is obtained by using the graph G_n of Theorem 3. The right inequality is obtained from the clique partition of K_n constructed in the proof of Theorem 4. It has x of its cliques in G and $\left(\frac{13}{30}n^2 - \frac{n}{6} - x\right)$ cliques in \overline{G} .

Concluding remarks

L. Pyber proved that the lower bound in Theorem 1 is sharp for n large. Possibly Theorem 3 is close to best possible; that is, max $\{cp(G)+cp(\overline{G})\}\sim 7n^2/25$ where the maximum is taken over all *n*-vertex graphs G. Suppose $G_1 \cup G_2 \cup G_3 = K_n$ where the G_i are edge-disjoint. If R is the graph diagrammed in Figure 3 with A= $=\overline{K}_{n/5}$, then we can have $G_1\cong G_2\cong R$ and so $cp(G_1)+cp(G_2)=2n^2/5$. (We use the same diagrammatic convention here as in Figure 1.) Probably this is the maximum possible value of $cp(G_1)+cp(G_2)$. The estimate $cc(G_1)+cc(G_2)+cc(G_3)=$ $=2n^2/5(1+o(1))$ was proved by L. Pyber (see pp. 393—398 of this issue). Perhaps

max $\{cc(G_1)+cc(G_2)+cc(G_3)\}=2n^2/5+5$, taking the maximum over all *n*-vertex graphs.

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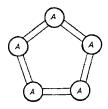


Fig. 3

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